

AMENDMENTS TO THE CLAIMS

1. (Original) A public key generation apparatus including:
 - a random number generator for generating a random number k_a that holds a relationship $0 < k_a < q$, where an element in a finite group F for which multiplication is defined is g and an order that is a prime number of the element g is q ; and
 - a public key generator for calculating a public key y_a in the finite group F from the random number k_a , the element g , and the prime number q ,
 - at least said random number generator and said public key generator being formed on one semiconductor integrated circuit, and
 - a controller of a first user as a distribution source of the public key controlling the random number generator and the public key generator for obtaining the public key y_a , and transmitting the obtained public key y_a to a second user as a distribution destination of the public key.
2. (Original) The public key generation apparatus of Claim 1 wherein
said public key generator calculates the public key y_a in the finite group F by a formula: $y_a = g^{k_a} \bmod q$, using the random number k_a , the element g , and the prime number q .
3. (Original) The public key generation apparatus of Claim 1 wherein
when the finite group F is an elliptic curve $E(F)$ in a finite field, and an element of the elliptic curve $E(F)$ is G ,
said public key generator calculates the public key y_a on the elliptic curve $E(F)$ by a formula: $y_a = k_a G \bmod q$, using the random number k_a , the element G , and the prime number q .
4. (Currently Amended) The public key generation apparatus of ~~any of Claims~~ Claim 1 to 3 wherein
said random number generator generates a new random number k_a after the calculation of the public key y_a is completed.

5. (Original) A shared key generation apparatus including:
 - a random number generator for generating a random number k_a that holds a relationship $0 < k_a < q$, where an element in a finite group F for which multiplication is defined is g and an order that is a prime number of the element g is q ; and
 - a shared key generator for calculating a shared key K_a in the finite group F from a public key y_b that is generated from a random number k_b which holds a relationship $0 < k_b < q$ and is generated by a second user as a distribution destination of the shared key, and the random number k_a ,
 - at least said random number generator and said shared key generator being formed on one semiconductor integrated circuit, and
 - a controller of a first user as a distribution source of the shared key obtaining the public key y_b from the second user as the shared key distribution destination, and controlling the random number generator and the shared key generator for deriving the shared key K_a .
6. (Original) The shared key generation apparatus of Claim 5 wherein
 - said shared key generator calculate the shared key K_a in the finite group F by a formula: $K_a = y_b^{k_a} \bmod q$, using the public key $y_b = g^{k_b} \bmod q$ which is generated by the second user as the shared key distribution destination and the random number k_a .
7. (Original) The shared key generation apparatus of Claim 5 wherein
 - when the finite group F is an elliptic curve $E(F)$ in a finite field and an element of the elliptic curve $E(F)$ is G ,
 - said shared key generator calculates the shared key K_a on the elliptic curve $E(F)$ by a formula: $K_a = k_a y_b \bmod q$, using the public key $y_b = k_b G \bmod q$ which is generated on the elliptic curve $E(F)$ from the random number k_b by the second user as the shared key distribution destination, and the random number k_a .
8. (Currently Amended) The shared key generation apparatus of ~~any of Claims~~ Claim 5 to 7 wherein

said random number generator generates a new random number k_a after the calculation of the shared key K_a is completed.

9. (Original) A key exchange apparatus including:

a random number generator for generating a random number k_a that holds a relationship $0 < k_a < q$, where an element in a finite group F for which multiplication is defined is g and an order that is a prime number of the element g is q ;

a public key generator for calculating a public key y_a in the finite group F from the random number k_a , the element g , and the prime number q ; and

a shared key generator for calculating a shared key K_a in the finite group F on the basis of the public key y_b generated from a random number k_b which holds a relationship $0 < k_b < q$ and is generated by a second user as a distribution destination of the shared key, and the random number k_a ,

at least said random number generator, said public key generator, and said shared key generator being formed on one semiconductor integrated circuit, and

a controller of a first user as a distribution source of the shared key controlling the random number generator and the public key generator for obtaining the public key y_b , and controlling the shared key generation unit for deriving the shared key k_a .

10. (Original) The key exchange apparatus of Claim 9 wherein

said public key generator calculates the public key y_a in the finite group F by a formula: $y_a = g^{k_a} \bmod q$, using the random number k_a , the element g , and the prime number q , and

said shared key generator calculates the shared key K_a in the finite group F by a formula: $K_a = y_b^{k_a} \bmod q$, using the public key $y_b = g^{k_b} \bmod q$ which is generated in the finite group F by the second user as the shared key distribution destination using the random number k_b , and the random number k_a .

11. (Original) The key exchange apparatus of Claim 9 wherein
when the finite group F is an elliptic curve E(F) in a finite field, and an element of the elliptic curve E(F) is G,
said public key generator calculates the public key y_a on the elliptic curve E(F) by a formula:
 $y_a = k_a G \bmod q$, using the random number k_a , the element G, and the prime number q, and
said shared key generator calculates the shared key K_a on the elliptic curve E(F) by a
formula: $K_a = k_a y_b \bmod q$, using the public key $y_b = k_b G \bmod q$ generated from the random number
 k_b on the elliptic curve E(F) by the second user as the shared key distribution destination, and the
random number k_a .

12. (Currently Amended) The key exchange apparatus of ~~any of Claims~~ Claim 9 to 11 wherein
the random number generator generates a new random number k_a after the calculation of the
public key y_a and the calculation of the shared key K_a are both completed.

13. (Original) A key exchange apparatus including:
a random number generator for generating a random number k_a that holds a relationship $0 < k_a < q$, where an element in a finite group F for which multiplication is defined is g and an order
that is a prime number of the element g is q;
a secret key holding unit for temporarily holding the random number k_a ;
a public key generator for calculating a public key y_a in the finite group F from the random
number k_a , the element g, and the prime number q; and
a shared key generator for calculating a shared key K_a in the finite group F using a public key
 y_b generated from a random number k_b which holds a relationship $0 < k_b < q$ and is generated by
a second user as a destination distribution of the shared key, and the random number k_a that is held
by the secret key holding unit,
at least said random number generator, said secret key holding unit, said public key generator,
and the shared key generator being formed on one semiconductor integrated circuit,

a controller of a first user as a distribution source of the shared key controlling the random number generator and the public key generator for obtaining the public key y_a , and transmitting the obtained public key y_a to a second user as a distribution destination of the shared key, and

said controller obtaining the public key y_b from the second user as the shared key distribution destination, and controlling the shared key generator for deriving the shared key K_a .

14. (Original) The key exchange apparatus of Claim 13 wherein

the public key generator calculates the public key y_a in the finite group F using the random number k_a , the element g , and the prime number q by a formula: $y_a = g^{k_a} \bmod q$, and

the shared key generator calculates the shared key K_a in the finite group F by a formula: $K_a = y_b^{k_a} \bmod q$, using the public key $y_b = g^{k_b} \bmod q$ that is generated in the finite group F from the random number k_b by the second user as the shared key distribution destination, and the random number k_a that is held in the secret key holding unit.

15. (Original) The key exchange apparatus of Claim 13 wherein

when the finite group F is an elliptic curve $E(F)$ in a finite field, and an element on the elliptic curve $E(F)$ is G ,

the public key generator calculates the public key y_a on the elliptic curve $E(F)$ using the random number k_a , the element G , and the prime number q by a formula: $y_a = k_a G \bmod q$, and

the shared key generator calculates the shared key K_a on the elliptic curve $E(F)$ by a formula: $K_a = K_{ayb} \bmod q$, using the public key $y_b = k_b G \bmod q$ that is generated from the random number k_b on the elliptic curve $E(F)$ by the second user as the shared key distribution destination, and the random number k_a that is held in the secret key holding unit.

16. (Currently Amended) The key exchange apparatus of ~~any of Claims~~ Claim 13 to 15 wherein

the random number generator generates a new random number k_a after the calculation of the public key y_a is completed, and

the secret key holding unit holds the new random number k_a generated by the random number generator.

17. (Currently Amended) The key exchange apparatus of ~~any of Claims~~ Claim 13 to 15 wherein the random number generator generates a new random number k_a after the calculation of the shared key K_a is completed, and

the secret key holding unit holds the new random number k_a generated by the random number generator.

18. (Currently Amended) A key exchanging method that employs the key exchange apparatus of ~~any of Claims~~ Claim 9 to 17, thereby exchanging the public keys that are generated by a first user and a second user that intend to exchange the public keys, respectively, to generate a shared key by the first user and the second user on the basis of the exchanged public key, respectively.

19. (New) A key exchanging method that employs the key exchange apparatus of Claim 13, thereby exchanging the public keys that are generated by a first user and a second user that intend to exchange the public keys, respectively, to generate a shared key by the first user and the second user on the basis of the exchanged public key, respectively.